Second Edition

Digital Electronics Drill

Solutions Test Drill I Test Drill II

A K Tripathi MD, Mechasoft Publishers and Educators and Asstt. Prof. Electronics Engineering Department IERT, Degree Division, Allahabad (INDIA)

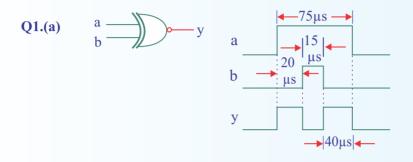
Mechasoft Publishers, Allahabad (INDIA)





Digital Electronics Drill

Test Drill I



Q2.
$$x^2 + 2x + 3 + (x - 2)^2 + 2(x - 2) + 1 = 3 \times 16 + 6 \Rightarrow 2x^2 = 50$$

 $x = \pm 5$

Choose positive integer value of x=5.

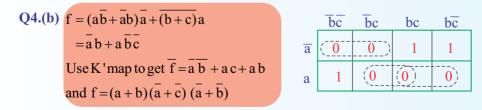
Q3.(a) The necessary and sufficient conditions for a function to be self dual are *(i)the function must be neutral .A function with equal number of max terms and min terms, is said to be neutral .*

(ii) the function must not contain mutually exclusive terms.

The term X = abc and the term Y = abc obtained by complementing each literal, are said to be mutually exclusive.

Only $f_1 = x$ and $f_4 = xy + \overline{xy} = y$ satisfy these conditions.

$$f_1(x, y, z) = x = \sum (4, 5, 6, 7)$$
 and $f_4(x, y, z) = y = \sum (2, 3, 6, 7)$

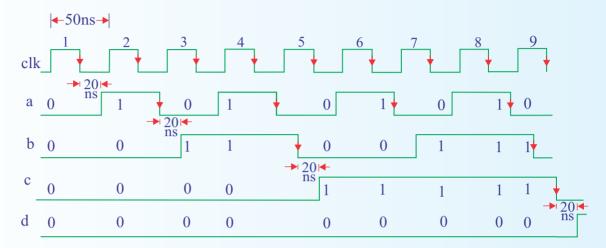


Q5.(d)
$$T_{CLK} = 1/(20 \times 10^6) = 0.05 \,\mu s = 50 \,ns$$

Let the counter state be defined as *dcba* with *d* as MSB.



As demonstrated below, the counter counts normally from dcba = 0000 to 0111 until clock 7. At the 8th clock 'a' transits from 0 to 1 before 'd' transits from 0 to 1. There by , the counter skips the state dcba=1000.



Q6.
$$1 \times R^{1} + 1 \times R^{0} + 2(R+1)^{1} + 3(R+1)^{0} = 1 \times R^{2} + 0 \times R^{1} + 2R^{0}$$
$$R^{2} - 3R - 4 = 0 \Longrightarrow R = -1 \text{ or } 4; R \text{ can not } \mathbf{be} \text{ negative}. R = 4$$

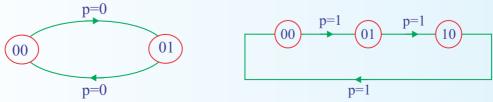
Q7.(d) $(26)_{16} \equiv (0010 \ 0110)_2$ $(36)_{16} \equiv (0011 \ 0110)_2$ $(A6)_{16} = (1010 \ 0110)_2$ and $(37)_8 = (011 \ 111)_2$ The only negative number in 2's complement is $(A6)_{16} = (1010 \ 0110)_2 \equiv -(90)_{10}$

Q8.(c) The output of NOR gate is

 $x = (a + b + c) + \overline{a} bc = \overline{a + b + c + bc} = \overline{a + b + c} = \overline{a} \overline{b} \overline{c}$

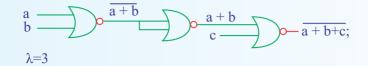
The input values a=b=c=0 will force the LED to glow.

Q9.(c) The machine is mod 2 counter for p=0 and mod 3 counter for p=1. It is self starting in the sense that it goes to valid counting sequence for any initial state.

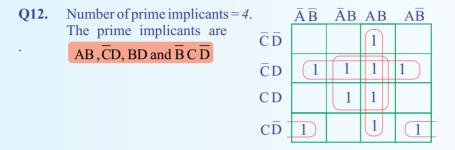








Q11.(b) Dividing a binary number by 6, is not possible by bit shifting.



Q13.(a) Inverting each bit and incrementing the result is equivalent to finding 2's complement. $(-124)_{10} \equiv (1,0000100)_2$ in 2's complement and 2's complement of $(1,0000100)_2$ is 0,1111100.

Q14.(d)

clock cycle	1	2	3	4	5	6	7	8	9	10
X	1	1	1	0	0	0	1	1	1	0
state	S ₀	S_1	S ₂	S ₃	S_0	S ₀	S ₀	\mathbf{S}_1	S ₂	S ₃
М	0	0	0	1	0	0	0	0	0	1

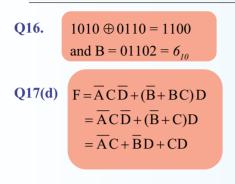
 \Rightarrow M=1 when input X has pattern '1110'

Q15. The clock period must be > (D FF delay + maximum combinational circuit delay + D FF setup time)

 $T_{CLK} > [15 + (35+10+25)+15]$ or $T_{CLK} > 100 \text{ ps}$ CLK frequency < $[100 \text{ ps}]^{-1} = 10 \text{ GHz}$

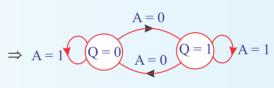
Solutions Test Drill I



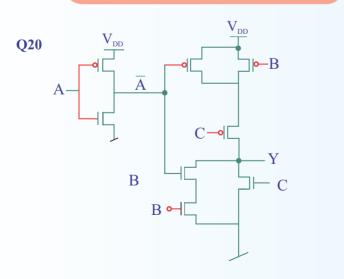


F	$\overline{C}\overline{D}$	ĒD	CD ;	CD
ĀB		1		Ĩ
ĀB			1	1
AB			1	
AB		[1		
		i		

Q18.(b)	Input	PS	FF input	NS
	А	Q	$D = out = \overline{A} D_0 + AD_1 = \overline{A} \overline{Q} + AQ$	Q^+
	0	0	1	1
	0	1	0	0
	1	0	0	0
	1	1	1	1



Q19.(b) $F = (\overline{A}\overline{D})C + (\overline{A}D)C + (\overline{A}\overline{D}).1 + ADB$ = $\overline{A}\overline{C}\overline{D} + \overline{A}CD + A\overline{D} + ABD$ Use K'map to get minimal SOP expression $F = \overline{A}CD + \overline{C}\overline{D} + AB + A\overline{D}$

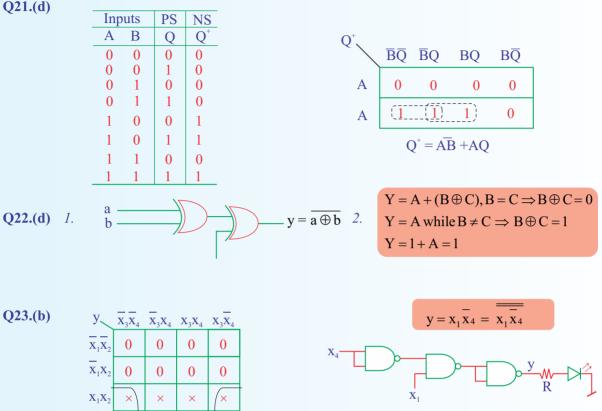


λ=4





Q21.(d)



- Q24.(c) The circuit is designed to do the subtraction operation, that is, X Y. For z = 1, $S_3S_2S_1S_0 = 0000$, that is, X - Y = 0 or X = Y.
- **Q25.(b)** Output frequency = f_{in} divided by 5×8×10=400.

0

×

×

 $\mathbf{X}_1 \mathbf{X}_2$ 1

$$\frac{J_{in}}{400} = 30 \times 10^3$$
 and $f_{in} = 400 \times 30 \times 10^3$ Hz = 12MHz



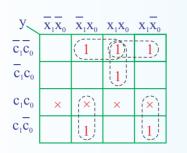


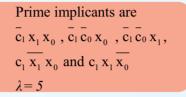


Digital Electronics Drill

Test Drill II







Q2.(b) For $t_{dlly}=0$ and considering only set up time constraints while ignoring hold time violations, the path from Q_1 to D_2 involves worst case delay =2× $t_{d,max}$ of INV+ $t_{cq,max}+t_{su}$ =2×1+1+2=5ns and there fore, the clock period $T_{CLK} \ge 5ns$. $f_{CLK,max}=(1/5)\times10^{-9}$ Hz=200 MHz.

The shortest delay from Q_2 to D_1 is $t_{eq,min}$ of FF2 + $t_{d,min}$ of INV = 0.5+0.5=1ns. There fore, there is hold time violation at input D_1 of FF1. While D_1 should be held constant for hold time, $t_h = 1.5$ ns, Q_2 should not be allowed to arrive at D_1 .

To eliminate this hold time violation at D_1 , it is required to delay Q_2 at least by 0.5ns so that Q_2 arrives at D_1 at least after 1.5 ns. This can be achieved by delaying CLK to FF2 by 0.5ns. Thus, $t_{dly,min} = 0.5ns$.

Q3.(b)

 8×1 MUX configuration gives

$$f = a + \overline{a} \, \overline{b} \, \overline{c} + \overline{a} \, \overline{b} \, \overline{c} = a + \overline{a} \, \overline{c} = a + \overline{c}$$

$$f_1 = \overline{a} + a\overline{c} = \overline{a} + \overline{c} \neq f$$

$$f_2 = \overline{a} \, \overline{b} \, \overline{c} + \overline{a} \, \overline{b} \, \overline{c} + a\overline{b} + a\overline{b} = \overline{a} \, \overline{c} + a(b + \overline{b}) = a + \overline{c} = f$$

$$f_3 = \overline{a} \, \overline{c} + \overline{a} \, \overline{c} \neq f$$

$$f_4 = \overline{b} \, \overline{c} + b\overline{c} + \overline{a} \, \overline{b} \, \overline{c} = \overline{b} \, \overline{c} + b(\overline{c} + \overline{a} \, \overline{c}) = \overline{b} \, \overline{c} + b(\overline{c} + \overline{a}) = (\overline{b} + b)\overline{c} + \overline{a} \, \overline{b} = \overline{a} \overline{b} + \overline{c} \neq \overline{z}$$

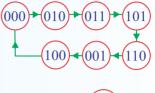




Q4.(d) 01000 11010₂=282₁₀ and $v_0 = (5/1025) \times 282 = 1.3756 V \neq 2.82V$.

Q5.(d)

PS	FF	Inputs		NS
$Q_1Q_2Q_3$	$D_1 = Q_3$	$D_2 = \overline{Q}_1 \overline{Q}_3 + Q_1 Q_3$	$D_3 = Q_2$	$Q_1^+Q_2^+Q_3^+$
1 1 1	1	1	1	1 1 1
0 0 0	0	1	0	0 1 0
0 1 0	0	1	1	0 1 1
0 1 1	1	0	1	1 0 1
1 0 1	1	1	0	1 1 0
1 1 0	0	0	: 1	0 0 1
0 0 1	1	0	0	1 0 0
1 0 0	0	0	0	0 0 0





The circuit is mod -7 counter without built in self-starting feature. From state $Q_1 Q_2 Q_3$ = 111, there is no ongoing.

	PS	Input	FF	Inp	out	NS	
	Q	Х	S=X	⊕Q	$R = \overline{X \oplus Q}$	Q^+	x=1
Q6.(b)	0	0	0		1	0	
	0	1	1		0	1	(Q=0) $(Q=1)$ $x=0$
	1	0	1		0	1	
	1	1	0		1	0	x=1

Q7. Let $v_0 = v_2$ for $b_3 b_2 b_1 b_0 = 0100$ and $v_0 = v_1$ for $b_3 b_2 b_1 b_0 = 0010$. Then, $v_2 = 2 v_1$ will ensure that circuit acts as a 4-bit DAC. For $b_3 b_2 b_1 b_0 = 0100$, $v_0 = v_2 = v_R \times (R_1/2R)$ For $b_3 b_2 b_1 b_0 = 0010$, $v_0 = v_1$ is calculated as follows.

$$\frac{-V_{R} - V}{R} = \frac{V}{2R} + \frac{V}{r} \Rightarrow V = \frac{-2rV_{R}}{2R + 3r}$$

$$\frac{0 - v_{1}}{R_{1}} = \frac{V}{r} \Rightarrow v_{1} = -\frac{R_{1}}{r} \left[\frac{-2rV_{R}}{2R + 3r} \right]$$

$$= \frac{2R_{1}V_{R}}{2R + 3r}$$

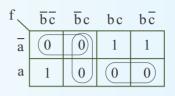
$$v_{2} = 2v_{1} \Rightarrow \frac{R_{1}V_{R}}{2R} = \frac{4R_{1}V_{R}}{2R + 3r} \Rightarrow \frac{r}{R} = k = 2$$

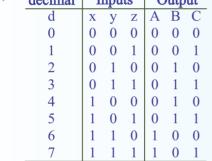


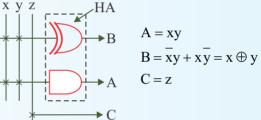


Step size, $S = 1.5/2^4 = 0.09375V$ and 1.0/0.09375 = 10.66. **Q8**. One step below, $(10)_{10} = (1010)_2$ and digital reading (1010)₂ represents $10 \times 0.09375 = 0.9375$ V. Error = 1.0 - 0.9375 = 0.0625 V. = 62.5 mV.

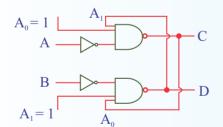
Q9.(b) $f(a, b, c) = \overline{S_0} (\overline{a} b + a \overline{b}) + S_0 \overline{(b + c)}$; $S_0 = a$ $=\overline{a}(\overline{a} \ b + \overline{a} \ \overline{b}) + \overline{a} \ \overline{b} \ \overline{c} = \overline{a} \ b + \overline{a} \ \overline{b} \ \overline{c}$ Use K' map to get $f' = \overline{ab} + \overline{b}c + ab$ $f = (f')' = \overline{ab} + \overline{b}c + ab$ =(a+b)(b+c)(a+b)Q10.(c) decimal | Inputs Output A B d х y Z С





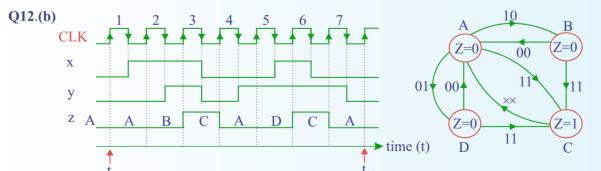


Q11.(c) The gate equivalent circuit is





S=A, R=B, Q=C and $\overline{Q} = D$



Note : The state diagram is the part of Q12. This has $n^2 t$ been printed in the book. Inconvenience is regretted.





Q13.(c) A table of transition from Q (present state) to Q^+ (next state) as per inputs J and K of JK flip flop together with required T input of T flip flop for corresponding state transition, is demonstrated below.

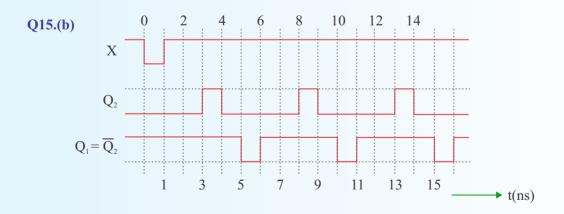
J	Κ	Q	Q^+	Т
0	0	0	0	0
0	0	1	1	0
0	1	0	0	0
0	1	1	0	1
1	0	0	1	1
1	0	1	1	0
1	1	0	1	1
1	1	1	0	1

 $T = \overline{J}KQ + J\overline{K}\overline{Q} + JK\overline{Q} + JKQ = (\overline{J}\overline{K}).0 + (\overline{J}K)Q + (J\overline{K})\overline{Q} + (JK).1$ Also, T = out = $\overline{J}\overline{K}.I_0 + \overline{J}K.I_1 + J\overline{K}.I_2 + JK.I_3$ Compare to get $I_0 = 0, I_1 = Q, I_2 = \overline{Q}$ and $I_3 = 1$

Q14.(d) The design of NOR gate demands 4 PMOS transistors in series and 4 NMOS transistors in parallel . The on resistance of a MOSFET is inversely proportional to W/L ratio . Each PMOS transistors should have W/L ratio equal to 4p so that the delay of NOR gate is equal to that of basic inverter .

The design of NAND gate demands 4 NMOS transistors in series and 4 PMOS transistors in parallel. In order that the delay of NAND gate is equal to that of basic inverter, the W/L ratio of each NMOS transistor should be 4n.

In the basic CMOS inverter, the W/L ratio 'p 'of PMOS transistor is about 2-3 times the W/L ratio 'n' of NMOS transistor so that the two transistors have equal on resistance. Therefore, the NOR gate requires the area much larger than the NAND gate.

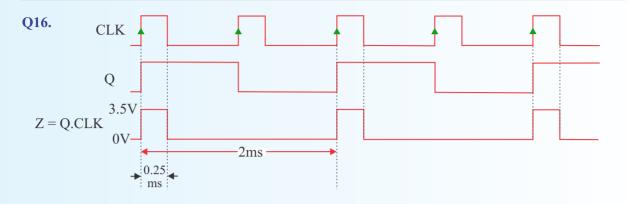


It is apparent from the waveforms sketched just above, that the outputs Q_1 and Q_2 are oscillatory in nature but the duty ratio at Q_1 is $(4/5) \times 100 = 80\%$ and that at Q_2 is $(1/5) \times 100 = 20\%$. The average power dissipated in resistor 5k Ω is

$$P_{av} = Peak Power \times (duty cycle at Q_1) = \frac{5^2}{5k\Omega} \times 0.8 = 4mW$$







The waveform at node Z has duty cycle = 0.25/2 = 0.125The average power dissipated in 360Ω resistor is $P_{av} = (3.5^2/360) \times 0.125$ W = 4.25 mW

Corrigendum:

In fig., the NAND gate should be AND gate . In answer key of Q16, read 4.25 instead of 8.5.Inconvenience is regretted

Q17.(c)
$$X = \overline{A}\overline{B}.0 + \overline{A}B.1 + A\overline{B}.0 + AB.A = AB + \overline{A}B = B$$

 $Y = \overline{C}X + C.C = B\overline{C} + C = B + C$

Q18.(c) The functionality is 4×1 MUX with b_1b_0 as select lines, a_0 to a_3 as data lines and f as output line.

 D_{0}/a_{0}		4×1		
 D_{1}/a_{1}		MUX		
 D_{2}/a_{2}		out	/1	
 D ₃ /a ₃	S_{1}/b_{1}	S ₀ /1	b_0	

Select	Output	
b ₁	b_0	f
0	0	a_0
0	1	a_1
1	0	a_2
1	1	a ₃

$$\phi(x_1, x_2, x_3) = x = x_1 x_2 + x_1 x_3$$

$$\phi(a, b, c) = y_1 = ab + ac$$

$$y_2 = \phi(y_1, a, b + c) = y_1 \overline{a} + \overline{y_1}(\overline{b + c})$$

$$= (a \overline{b} + \overline{a} \overline{c})\overline{a} + \overline{ab + ac}.\overline{b}c = \overline{ac} + \overline{ab}.\overline{ac}.\overline{b}c$$

$$= \overline{ac} + (\overline{a} + b)(a + c)\overline{b}c = \overline{ac} + [\overline{a}c + ab + bc]\overline{b}c = \overline{ac} + \overline{ab}c$$

$$y = \overline{y_2}.\overline{b} + \overline{c} = \overline{y_2}.\overline{b}c = (\overline{ac} + \overline{ab}c)\overline{b}c = \overline{abc} = a + b + \overline{c}$$

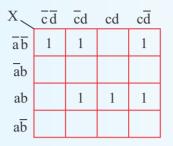
_

x	$\overline{\mathbf{X}}_{2}\overline{\mathbf{X}}_{3}$	$\overline{\mathbf{X}}_{2}\mathbf{X}_{3}$	X_2X_3	$\mathbf{X}_{2}\mathbf{\overline{X}}_{3}$		
$\overline{\mathbf{x}}_{1}$	0	1	1)	0		
Х	1	1)	0	0		
$\mathbf{x} = \mathbf{x}_1 \overline{\mathbf{x}}_2 + \overline{\mathbf{x}}_1 \mathbf{x}_3$						





Q20.(d) $X = \overline{abc} \ \overline{abd} \ \overline{abd} \ \overline{abd} \ \overline{abc} = \overline{abc} + \overline{abd} + abd + abc$



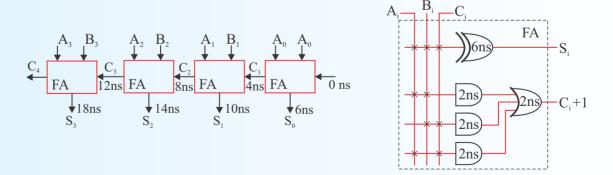
The LED turns on when X = 0 and this is true when input is greater than 2 but less than 13.

Decimal	а	b	с	d	Х
0	0	0	0	0	1
1	0	0	0	1	1
2 3	0	0	1	0	1
3	0	0	1	1	0
4	0	1	0	0	0
5	0	1	0	1	0
6	0	1	1	0	0
7	0	1	1	1	0
8	1	0	0	0	0
9	1	0	0	1	0
10	1	0	1	0	0
11	1	0	1	1	0
12	1	1	0	0	0
13	1	1	0	1	1
14	1	1	1	0	1
15	1	1	1	1	1

Q21.(b) $X = (a \oplus b)(\overline{b \oplus c})c = (\overline{a}b + a\overline{b})(\overline{bc} + bc)c$ = $(\overline{a}b + a\overline{b})bc = \overline{a}bc$

For abc = 011, the transistor terms on and so only the LED.

Q22. As demonstrated below, total time for a successful addition operation = 18ns.



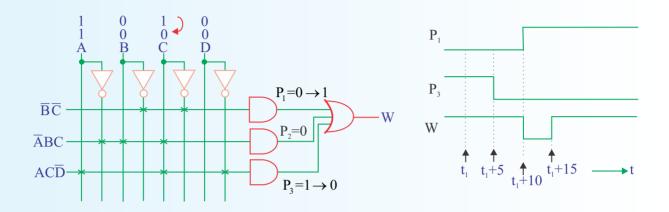
Q23.(d) $f = \overline{Z}D + ZB = \overline{Z}(\overline{X}Y + X\overline{Y}) + Z\overline{X}Y$ For Z = 1, f = 0 indicates either X = 0 and Y = 0 or X = 1 and Y = 0/1.

Q24.(d) In fact, the input x = 0111 is correctly matched to output y = 0011.





Q25.(b) The input C changes from 1 to 0 while inputs A, B and D remain unchanged.







MECHASOFT PUBLISHESRS accelerating pace of learning



Publishers and Educators